

Solution 9: Data structures

ETH Zurich

1 Choosing data structures

1. You can use a doubly-linked list. An arrayed list is also suitable if it is implemented as a circular buffer (that is, the list can start from any element in the array), in which case inserting in the beginning of the list is also efficient. A disadvantage of an arrayed list is that adding a station will sometimes take longer (when the array does not have any more free slots and has to be reallocated), an advantage is fast access by index, which is not mentioned in the scenario, but is always good to have.

A disadvantage of a doubly-linked list is high memory overhead: in addition to the reference to a station object each list element stores two other references (to the next and the previous element). Arrayed list also has a memory overhead (free array slots), however for common implementations this overhead will not be as high.

2. A hash table with names (strings) as keys and phone numbers as values, because hash table allows efficient access by key.
3. A stack, because the step that was added last is always the first to roll back.
4. A linked list, because it supports efficient insertion of the elements of the second list into the proper place inside the first list while merging. The insertion is done by re-linking existing cells and does not require creating a copy of either of the lists.
5. A queue, because the first call added to the data structure should be the first one to be processed.

2 Short trips: take two

Listing 1: Class *SHORT_TRIPS*

```
note
  description: "Short trips."

class
  SHORT_TRIPS

inherit
  ZURICH_OBJECTS

feature -- Explore Zurich

  highlight_short_distance (s: STATION)
    -- Highlight stations reachable from 's' within 3 minutes.
  require
```

```
    station_exists: s /= Void
do
  create times
  highlight_reachable (s, 3 * 60)
end

feature {NONE} -- Implementation

times: V_HASH_TABLE [STATION, REAL_64]
  -- Table that maps a station to the maximum time that was left after visiting that
  -- station.
  -- Stations that were never visited, are not in the table.

highlight_reachable (s: STATION; t: REAL_64)
  -- Highlight stations reachable from 's' within 't' seconds.
require
  station_exists: s /= Void
local
  line: LINE
  next: STATION
do
  if t >= 0.0 and (not times.has_key (s) or else times [s] < t) then
    times [s] := t
    Zurich_map.station_view (s).highlight
  across
    s.lines as li
  loop
    line := li.item
    next := line.next_station (s, line.north_terminal)
    if next /= Void then
      highlight_reachable (next, t - s.position.distance (next.position) / line.speed)
    end
    next := line.next_station (s, line.south_terminal)
    if next /= Void then
      highlight_reachable (next, t - s.position.distance (next.position) / line.speed)
    end
  end
end
end
end
end
```

3 English to Swiss-German Dictionary

Listing 2: Class *TRIE*

```
note
  description: "Trie data structure, where child nodes are stored in a linked list,
  sorted by label."

class
  TRIE
```

feature -- Access

```
item (k: STRING): STRING
  -- Value associated with key 'k'.
require
  key_not_void: k /= Void
local
  e: EDGE
do
  if k.is_empty then
    Result := value
  else
    from
      e := first
    until
      e = Void or else e.label >= k [1]
    loop
      e := e.next
    end
  if e /= Void and then e.label = k [1] then
    Result := e.node.item (k.substring (2, k.count))
  end
end
end

to_string: STRING
  -- String representation with keys sorted lexicographically.
do
  Result := to_string_with_prefix ("")
ensure
  result_not_void: Result /= Void
end
```

feature -- Element change

```
insert (k, v: STRING)
  -- Insert key–value pair (k, v).
require
  key_not_void: k /= Void
local
  e, found: EDGE
do
  if k.is_empty then
    value := v
  else
    if first = Void or else first.label > k [1] then
      create found.make (k [1], create { TRIE })
      found.set_next (first)
      first := found
    else
      from
```

```

        e := first
    until
        e.next = Void or else e.next.label > k [1]
    loop
        e := e.next
    end
    if e.label = k [1] then
        found := e
    else
        create found.make (k [1], create {TRIE})
        found.set_next (e.next)
        e.set_next (found)
    end
end
found.node.insert (k.substring (2, k.count), v)
end
ensure
    correct_value: item (k) ~ v
end

```

feature {TRIE} *-- Implementation*

first: EDGE

-- Edge connecting to the first child node.

value: STRING

-- Value stored in the node.

-- 'Void' if the node doesn't correspond to a whole key.

to_string_with_prefix (key_prefix: STRING): STRING

-- String representation with keys sorted lexicographically

-- and 'key_prefix' prepended to all keys.

require

key_prefix_not_void: key_prefix /= Void

local

e: EDGE

do

if *value /= Void* **then**

Result := *key_prefix + " - " + value + "%N"*

else

Result := ""

end

from

e := first

until

e = Void

loop

Result.append (*e.node.to_string_with_prefix (key_prefix + e.label.out)*)

e := e.next

end

ensure

result_not_void: Result /= Void

```
end  
end
```