## **Introduction to Eiffel**

 $\odot$ 

1

Martin Nordio ETH Zurich martin.nordio@inf.ethz.ch

Distributed and Outsourced Software Engineering - ETH course, Fall 2012

# CONTRACTS

### Contracts

A contract is a semantic condition characterizing usage properties of a class or a feature

Three principal kinds:

- Precondition
- Postcondition
- Class invariant

## **Design by Contract**

Together with the implementation ("*how*") of each software element, describe "*what*" it is supposed to do: its contract

Three basic questions about every software element:

What does it assume?

> What does it guarantee?
> What does it maintain?
Postcondition
Invariant

## **Contracts in programming languages**

Eiffel: integrated in the language

Java: Java Modeling Language (JML), iContract etc.

.Net languages: Code Contracts (a library)

Spec# (Microsoft Research extension of C#): integrated in the language

UML: Object Constraint Language

Python

C++: Nana

### **Precondition**

Property that a feature imposes on every client:



A client calling a feature must make sure that the precondition holds before the call

A client that calls a feature without satisfying its precondition is faulty (buggy) software.

### Another example:

```
extend (a_element: G)

require

valid_elem: a_element /= void

not_full: not is_full

do ... end
```

```
A feature with a require clause

require

label_1: cond_1

label_2: cond_2 ...

label_n: cond_n

is equivalent to

require
```

label: cond\_1 and cond\_2 and ...

cond\_n





### Postconditions

Precondition: obligation for clients Postcondition: benefit for clients

```
extend (a_element: G)
ensure
inserted: i_th (count) = a_element
```

```
index (a_element: G): INTEGER
```

#### ensure

```
exists: result > 0 implies i_th (result) = a_element
no_exists: result = -1 implies not is_inserted (a_element)
```

### **Old notation**

Usable in postconditions only

Denotes value of an expression as it was on routine entry

Example (in a class ACCOUNT):

```
balance: INTEGER
-- Current balance.
```

```
deposit (v: INTEGER)
    -- Add v to account.
    require
        positive: v > 0
        do
        ...
    ensure
        added: balance = old balance + v
    end
```

A feature must make sure that, if its precondition held at the beginning of its execution, its postcondition will hold at the end.

A feature that fails to ensure its postcondition is buggy software.

### class

BANK\_ACCOUNT create make feature make (n : STRING) -- Set up with name n

> require n /= Void do

> > name := n balance := 0

ensure name = n end

name : STRING balance : INTEGER deposit (v: INTEGER) -- Add amount v do balance := balance + vensure balance = old balance + vend invariant name /= Void balance >= 0

end

Issues: what happens, under inheritance:

Invariant Inheritance rule:

The invariant of a class automatically includes the invariant clauses from all its parents,

"and"-ed.

When redeclaring a routine, we may only:

Keep or weaken the precondition

Keep or strengthen the postcondition

A simple language rule does the trick!

Redefined version may have nothing (assertions kept by default), or

require else new\_pre
ensure then new\_post

Resulting assertions are:

original\_precondition or new\_pre

original\_postcondition and new\_post