# Einführung in die Programmierung Introduction to Programming 

Prof. Dr. Bertrand Meyer

## Exercise Session 4

## Problems in Assignment-2 Solutions

## > Command or query?

- connecting_lines
(a_station_1, a_station_2: STATION): V_SEQUENCE [LINE]
- Noun phrases for query names; verb phrases for command names
> Instruction separation?
- Comma (,), space( ), semicolon (;), or nothing
> STRING_8 Vs. STRING_32
make

```
local
    l_line: STRING_32
            c: UTF_CONVERTER
```

do
Io.read_line
l_line := c.utf_8_string_8_to_string_32 (Io.last_string)
print (l_line.count)
end

## Today

> Understanding contracts
(preconditions, postconditions, and class invariants)
> Reference types vs. expanded types
> Basic types
> Entities and objects
Object creation
Assignment

## Why do we need contracts?

- They are executable specifications that evolve together with the code
- Together with tests, they are a great tool for finding bugs
- They help us reason about an O-O program at the level of classes and routines
- Proving (part of) programs correct requires some way to specify how the program should operate. Contracts are a way to specify the program


## Assertions



## Assertion clause

The assertion tag (if present) is used to construct a more informative error message when the condition is violated.

Property that a feature imposes on clients
clap ( $n$ : INTEGER)
-- Clap $n$ times and update count.

```
    require
    not_too_tired: count <= 10
    n_positive: n>0
```

A feature without a require clause is always applicable, as if the precondition reads
require always_OK: True

## Postcondition

Property that a feature guarantees on termination

```
clap (n: INTEGER)
                            -- Clap n times and update count.
        require
        not_too_tired: count <= 10
        n_positive: n>0
    ensure
        count_updated: count = old count + n
```

A feature without an ensure clause always satisfies its postcondition, as if the postcondition reads ensure
always_OK: True

## Class Invariant

Property that is true of the current object at any observable point
class ACROBAT
invariant
count_non_negative: count $>=0$
end

A class without an invariant clause has a trivial invariant
always_OK: True

Add pre- and postconditions to:
smallest_power (n, bound: NATURAL): NATURAL
-- Smallest $x$ such that ${ }^{\prime} n^{\wedge} x$ is greater or equal ` bound'.
require
???
do
ensure
???
end

## One possible solution

smallest_power ( $n$, bound: NATURAL): NATURAL
-- Smallest $x$ such that ${ }^{n} n^{\wedge} x$ is greater or equal `bound'. require
n_large_enough: $n>1$
bound_large_enough: bound >1 do

## ensure

greater_equal_bound: $\mathrm{n}^{\wedge}$ Result >= bound smallest: $\mathrm{n}^{\wedge}$ (Result - 1) < bound end

## Hands-on exercise

Add invariants) to the class ACROBAT_WITH_BUDDY.

Add preconditions and postconditions to feature make in ACROBAT_WITH_BUDDY.

## Class ACROBAT_WITH_BUDDY

```
class
    ACROBAT_WITH_BUDDY
inherit
    ACROBAT
        redefine
            twirl, clap, count
        end
create
    make
feature
make (p: ACROBAT)
            do
        -- Remember ` p' being
        -- the buddy.
    end
```

```
clap (n: INTEGER)
        do
            -- Clap `n' times and
            -- forward to buddy.
        end
```

    twirl (n: INTEGER)
        do
            -- Twirl 'n' times and
            -- forward to buddy.
        end
    count: INTEGER
    do
        -- Ask buddy and return his
            -- answer.
        end
    buddy: ACROBAT
    end

## What are reference and expanded types?

Reference types: s contains the address (reference or location) of the object.
Example:
s: STATION


Expanded types: $p$ points directly to the object.
Example:

$$
p: \text { POINT }
$$



## Why expanded types?

To represent basic types (INTEGER, REAL,...)

To model external world objects realistically, i.e. objects that have sub-objects (and no sharing), for example a class WORKSTATION and its CPU.

## How to declare an expanded type?

To create an expanded type, declare the class with keyword expanded:
expanded class COUPLE
feature -- Access
man, woman: HUMAN
years_together: INTEGER
end

Now all the entities of type COUPLE are automatically expanded:
pitt_and_jolie: COUPLE
Expanded

## Objects of reference or expanded types

Objects of reference types: they don't exist when we declare them (they are initially Void).

$$
s: \text { STATION }
$$

We need to explicitly create them with a create instruction.
creates
Objects of expanded types: they exist by just declaring them (they are never Void)
$p:$ POINT
Feature default_create from ANY is implicitly invoked on them

## Can expanded types contain reference types?

Expanded types can contain reference types, and vice versa.



$$
a=b ?
$$


(VECTOR)


$$
a=b ?
$$

## Expanded entities equality

## 

$$
a=b ?
$$

Entities of expanded types are compared by value!

## Expanded entities equality



$$
a=b ?
$$

## Expanded entities equality



$$
a=b ?
$$

## Basic types

Their only privilege is to use manifest constants to construct their instances:
b: BOOLEAN
x. INTEGER
c: CHARACTER
s: STRING
$b:=$ True
$x:=5 \quad--$ instead of create $x . m a k e \_f i v e$
$c:=$ ' $c$ '
$s:=$ "I love Eiffel"

## Basic types

Some basic types (BOOLEAN, INTEGER, NATURAL, REAL, CHARACTER) are expanded...
$a:=b$

... and immutable (they do not contain commands to change the state of their instances)...


## Strings are a bit different

Strings in Eiffel are not expanded...
s: STRING

... and not immutable
$s:=$ "I love Eiffel"
s.append (" very much!")

## Object comparison: = versus $\sim$

```
s1: STRING = "Teddy"
s2: STRING = "Teddy"
si = s2 -- False: reference comparison on different objects
```

s1~s2 --True

Now you know how to compare the content of two objects

## Initialization

Default value of any reference type is Void Default values of basic expanded types are:
> False for BOOLEAN
> 0 for numeric types (INTEGER, NATURAL, REAL)
> "null" character (its code is 0 ) for CHARACTER Default value of a non-basic expanded type is an object, whose fields have default values of their types


## Initialization

What is the default value for the following classes?
expanded class POINT feature $x, y$. REAL end

class VECTOR
feature $x, y$. REAL end
Void

STRING
Void

## Creation procedures

$>$ Instruction create $x$ will initialize all the fields of the new object attached to $x$ with default values
$>$ What if we want some specific initialization? E.g., to make object consistent with its class invariant?

Class CUSTOMER
id: STRING
invariant

$$
i d /=\text { Void }
$$


> Use creation procedure:
create a_customer.set_id ("13400002")

## Class CUSTOMER

class CUSTOMER

-- Unique identifier for Current.
set_id (a_id: STRING)

require
id_exists : a_id/= Void
do
$i d:=a \_i d$
ensure
id_set : id = a_id
end

| invariant <br> id_exists : id $/=$ Void |
| :---: |
| end |

## Object creation

To create an object:

- If class has no create clause, use basic form: create $x$
> If the class has a create clause listing one or more procedures, use
create x.make (...)
where make is one of the creation procedures, and (...) stands for arguments if any.


## Some acrobatics



```
class DIRECTOR
create prepare_and_play
feature
    acrobat1, acrobat2, acrobat3: ACROBAT
    friend1, friend2: ACROBAT_WITH_BUDDY
    author1: AUTHOR
    curmudgeon1: CURMUDGEON
    prepare_and_play
    do
        author1.clap (4)
        friend1.twirl (2)
        curmudgeon1.clap (7)
        acrobat2.clap (curmudgeon1.count)
        acrobat3.twirl (friend2.count)
        friend1.buddy.clap (friend1.count)
        friend2.clap (2)
    end
end
```


## Some acrobatics

prepare_and_play
do
create acrobat1
create acrobat2
create acrobat3

```
class DIRECTOR
class DIRECTOR
create prepare_and_play
create prepare_and_play
feature
feature
    acrobat1, acrobat2, acrobat3: ACROBAT
    acrobat1, acrobat2, acrobat3: ACROBAT
    friend1, friend2: ACROBAT_WITH_BUDDY
    friend1, friend2: ACROBAT_WITH_BUDDY
    author1: AUTHOR
    author1: AUTHOR
    curmudgeon1: CURMUDGEON
    curmudgeon1: CURMUDGEON
    prepare_and_play
    prepare_and_play
        do
        do
1 create acrobat1
1 create acrobat1
2 create acrobat2
2 create acrobat2
3 create acrobat3
3 create acrobat3
4 create friend1.make_with_buddy (acrobat1)
4 create friend1.make_with_buddy (acrobat1)
5 create friend2.make_with_buddy (friend1)
5 create friend2.make_with_buddy (friend1)
6 create author1
6 create author1
7 create curmudgeon1
7 create curmudgeon1
    end
    end
end

Which of the classes mentioned here have creation procedures?

Which entities are still Void after execution of line 4?

\section*{Custom initialization for expanded types}
\(>\) Expanded classes are not creatable using a creation feature of your choice
expanded class POINT
create make
feature make do \(x:=5.0 ; y:=5.0\) end
end
>But you can use (and possibly redefine) default_create expanded class POINT
inherit ANY
redefine default_create
feature
defaul__create
do
end \(^{x:=5.0 ; y:=} 5.0\)
end

\section*{Assignment}
>Assignment is an instruction (What other instructions do you know?)
>Syntax:
\[
a:=b
\]
> where \(a\) is a variable (e.g., an attribute) and \(b\) is an expression (e.g. an argument or a query);
\(>a\) is called the target of the assignment and \(b\) the source.
\(>\) Semantics:
> after the assignment \(a\) equals \(b(a=b)\);
> the value of \(b\) is not changed by the assignment.

\section*{Reference assignment}

a references the same object as \(b\) :
\[
a=b
\]

\section*{Expanded assignment}

\[
a:=b
\]

The value of \(b\) is copied to \(a\), but again:
\[
a=b
\]

\section*{Assignment}

Explain graphically the effect of an assignment:

\[
a:=b
\]

Here COUPLE is an expanded class, HUMAN is a reference class

\section*{Attachment}

More general term than assignment
> Includes:
> Assignment
\[
a:=b
\]
> Passing arguments to a routine
\(f\) (a: SOME_TYPE) do ... end
\(f(b)\)
> Same semantics

\section*{Dynamic aliasing}
\(a, b:\) VECTOR
create b.make (1.0, 0.0)
\(a:=b\)

> now \(a\) and \(b\) reference the same object (they are two names or aliases of the same object)
\(>\) any change to the object attached to a will be reflected when accessing it using \(b\)
\(>\) any change to the object attached to \(b\) will be reflected when accessing it using a

\section*{Dynamic aliasing}

What are the values of \(a . x, a . y, b . x\) and b.y after executing instructions 1-4?
\(a, b:\) VECTOR
create a.make (-1.0, 2.0)


4 a.set_y (-10.0)

Meet Teddy
```

