Mock Exam 2

ETH Zurich

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Name: ____

Group: _____

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1 Contracts (11 points)

We are interested in a software system simulating a cellular automaton. The universe is represented by a finite square grid composed of square cells (there is at least 1). Each cell can be in two states: alive or dead. Every cell, depending on its position in the grid, can have from a minimum of 3 neighbors (a cell in a corner) to a maximum of 8 neighbors (a cell in the middle).

The evolution of the automaton from one generation to the next is fully determined by the following set of rules:

- Any living cell with less than 2 living neighbors dies in the next generation.
- Any living cell with 2 or 3 living neighbors lives in the next generation.
- Any living cell with more than 3 living neighbors dies in the next generation.
- Any dead cell with exactly 3 living neighbors becomes alive in the next generation.
- Any dead cell with a number of living neighbors different from 3 stays dead in the next generation.

The evolution from one generation into the next happens by applying the above rules simultaneously to every cell in the grid (see Figures 1 and 2).



Figure 1: Sample first generation. A black square is a living cell.

Figure 2: Second generation, computed from the first according to the given set of rules.

Your task is to add appropriate contracts (preconditions, postconditions and class invariants) to the excerpt of class $CELL_GRID$ below, so that the informal specification above and the feature comments are reflected in each class interface.

Please note that the number of dotted lines does not indicate the number of missing contracts. It might also be useful to have a look at the excerpt of class $ARRAY_2$ shown below.

1.1 Solution

```
class

CELL_GRID

create

make
```

feature {*NONE*} -- Initialization

```
make (a_dimension: INTEGER)
            -- Initialize grid's dimension to 'a_dimension' and its cells to dead.
       require
            dim_positive: a_dimension >= 1
       do
            -- Implementation omitted.
       ensure
           dim\_set: dim = a\_dimension
             current_grid_initialized_to_default : current_grid. all_default
       end
feature -- Access
    dim: INTEGER
            -- Grid dimension.
    cell_at (i, j: INTEGER): BOOLEAN
            -- Value of cell at (i,j).
       require
            i\_within\_bounds: i >= 1 \text{ and } i <= dim
           j_within_bounds: j \ge 1 and j \le dim
       do
            -- Implementation omitted.
       ensure
            right\_cell : Result = current\_grid.item (i, j)
       end
 feature -- Status Setting
    set_cell_status (b: BOOLEAN; i, j: INTEGER)
            -- Set status of cell at (i,j).
       require
            i_within_bounds: i \ge 1 and i \le dim
            j_within_bounds: j \ge 1 and j \le dim
       do
            -- Implementation omitted.
       ensure
            cell_status_set : cell_at (i, j) = b
       end
feature -- Basic operations
    compute\_next\_generation
```

-- Compute next_grid, copy it to current_grid and re–initialize <code>next_grid</code> .

-- Implementation omitted

end

do

feature {*NONE*} -- Implementation

```
current_grid: ARRAY2 [BOOLEAN]
```

```
-- Grid representation as a matrix of boolean cells ("True" means alive for a
           cell).
new_state_of_cell (i, j, living_neighbors: INTEGER): BOOLEAN
       -- Apply Conway's Game of Life rules to compute new state for cell at (i,j)
           given a number of 'living_neighbors'.
   require
       i_within_bounds: i \ge 1 and i \le dim
       j_within_bounds: j \ge 1 and j \le dim
       living_neighbors_within_bounds: living_neighbors >= 0 and living_neighbors <=
           8
   do
      -- Implementation omitted.
   ensure
       death\_rule\_1: current\_grid.item (i, j) and (living\_neighbors < 2 \text{ or})
           living_neighbors > 3) implies not Result
        life_rule : current_grid.item (i, j) and (living_neighbors = 2 \text{ or})
           living_neighbors = 3) implies Result
       birth_rule: not current_grid.item(i, j) and (living_neighbors = 3) implies
           Result
       death_rule_2: not current_grid.item (i, j) and (living_neighbors = 3) implies
           not Result
   end
```

invariant

```
current_grid_exists : current_grid /= Void
grid_dimension_positive : dim > 0
current_grid_dimension_is_dim: current_grid.width = dim and current_grid.height = dim
end
```

2 Data Structures (16 points)

In this task you are going to implement several operations for a generic class SET[G].

A set is a collection of distinct objects. Every element of a set must be unique; no two members may be identical. All set operations preserve this property. The order in which the elements of a set are listed is irrelevant (unlike for a sequence or tuple). Therefore the two sets $\{5, 10, 12\}$ and $\{10, 12, 5\}$ are identical.

There are several fundamental operations for constructing new sets from given sets.

- Union: The union of A and B, denoted by $A \cup B$, is the set of all elements that are members of either A or B.
- Intersection: The intersection of A and B, denoted by $A \cap B$, is the set of all elements that are members of both A and B.
- Relative complement of B in A (also called the set-theoretic difference of A and B), denoted by $A \setminus B$ (or A B), is the set of all elements that are members of A but not members of B.

The Jaccard index (or coefficient) measures similarity between sample sets, and is defined as the size of the intersection divided by the size of the union of the sample sets (see Figure 3). If both sets are empty the Jaccard coefficient is defined as 1.0.

$$J(A,B) = \frac{|A \cap B|}{|A \cup B|}$$

Figure 3: Jaccard index definition for non-empty sets A and B.

Your task is to fill in the gaps of class SET[G] below. Please note:

- Your code should satisfy the contracts and provide new contracts where necessary.
- The set should never contain **Void** elements.
- The number of dotted lines does not indicate the number of missing contract clauses or code instructions.
- The implementation of class SET[G] is based on an arrayed list. The arrayed list is set up to use object comparison, so features like *has* and *prune* use object equality instead of reference equality when comparing elements from the set. The following features of class $ARRAYED_LIST$ may be useful:

class $ARRAYED_LIST$ [G]

feature

has (v: G): BOOLEAN -- Does current include 'v'?

start

-- Move cursor to first position if any.

extend (v: G) -- Add 'v' to the end. prune (v: G)
-- Remove first occurrence of 'v', if any, after cursor position.
-- Move cursor to right neighbor.
-- Other features are omitted.
end

2.1 Solution

```
class
  SET\left[ \mathbf{G} \right]
create
  make_empty
feature {NONE} -- Initialization
  make_empty
      -- Create empty Current.
    do
     create content.make (0)
      content.compare_objects
    ensure
      empty_content: content.is_empty
    end
feature -- Access
  count: INTEGER
      -- Cardinality of the current set.
    do
      Result := content.count
    end
  is_empty: BOOLEAN
      -- Is current set empty?
    do
      Result := count = 0
    end
  has (v: G): BOOLEAN
      -- Does current set contain 'v'?
    require
      v \mid = Void
    do
     Result := content.has(v)
    end
  add (v: G)
      -- Add 'v' to the current set.
   require
```

```
v \mid = Void
    do
      if not has (v) then
        content. extend (v)
      end
    ensure
      in\_set\_already: old has(v) implies (count = old \ count)
      added_{to\_set}: not old has(v) implies (count = old \ count + 1)
    end
  remove (v: G)
      -- Remove 'v' from the current set.
    require
      v \mid = Void
    do
      if has (v) then
        content. start
        content. prune(v)
      end
    ensure
      removed\_count\_change: old has(v) implies (count = old \ count - 1)
      not\_removed\_no\_count\_change: not old has(v) implies (count = old count)
      item_deleted: not has (v)
    end
  duplicate: like Current
      -- Deep copy of Current.
    do
      create Result.make_empty
      across content as c
      loop
        Result.add (c.item)
      end
    ensure
      same\_size: Result.count = count
      same_content: across content as c all Result.has (c.item) end
    end
feature -- Set operations.
  union (another: like Current): like Current
      -- Union product of the current set and 'another' set.
    require
      another = Void
    do
      Result := another.duplicate
      across content as c
      loop
        Result.add (c.item)
      end
    ensure
      not\_smaller: Result.count >= count and Result.count >= another.count
```

end

```
intersection (another: like Current): like Current
      -- Intersection product of the current set and 'another' set.
   require
      another = Void
   do
     create Result.make_empty
      across content as c
     loop
        if another.has (c.item) then
         Result.add (c.item)
       end
     end
   ensure
      not_{bigger}: Result.count <= count and Result.count <= another.count
   end
  difference (another: like Current): like Current
      -- Set-theoretic difference of the current set and 'another' set.
   require
      another = Void
   do
     create Result.make_empty
      across content as c
     loop
        if not another.has (c.item) then
         Result.add (c.item)
       end
     end
   ensure
      not\_bigger\_than: Result.count <= count
      not\_smaller\_than: Result.count >= count - another.count
   end
feature -- Set metrics.
  jaccard_index (another: like Current): REAL_64
      -- Jaccard similarity coefficient between current set and 'another' set.
   require
      another = Void
   do
      if not (is_empty and another.is_empty) then
       Result := intersection (another).count / union (another).count
     else
       Result := 1.0
     end
    ensure
      bounds: Result \geq 0.0 and Result \leq 1.0
      empty_case: (is_empty and another.is_empty) implies \mathbf{Result} = 1.0
   end
```

feature {*NONE*} -- Implementation

content: $ARRAYED_LIST[G]$ -- Items of the set.

invariant

 $content_exists: content /=$ Void $content_object_comparison: content.object_comparison$ $non_negative_cardinality: count >= 0$

 \mathbf{end}

3 Recursion (14 points)

The N-queens problem is the problem of positioning N queens on an $N \times N$ board such that no queen can attack another (i.e., share the same row, column, or diagonal). The N-queens problem can be solved recursively: having a solution for the first 4 rows of the board can be used to build a solution for the 5th row, as is being done in Figure 4.



Figure 4: An example of a partial solution

A safe location is one which cannot be attacked by any of the currently placed queens.

A routine to solve the N-queens problem, *complete* (*partial*: *SOLUTION*), does as follows: if the partial solution is not yet complete, then for each safe location in the current row, add the safe location to the solution and use this new solution to solve the problem for the next row. The *current row* is *partial*.*row_count* + 1; for example in Figure 4 the partial solution has row_count equal to 4, thus the current row is 5. If the solution is already complete then it is added to the list of solutions.

You must complete the implementation of *PUZZLE* (which has an attribute *solutions* to store all solutions) below by filling in the body of *complete* and *attack_each_other*. Note that a solution can be added to the list of solutions using the *extend* feature from *LIST*.

3.1 Solution

note description: "N-queens puzzle." class PUZZLE feature -- Access size: INTEGER -- Size of the board. solutions: LIST [SOLUTION] -- All solutions found by the last call to 'solve'. feature -- Basic operations solve (n: INTEGER) -- Solve the puzzle for 'n' queens.

```
require
      solvable: n > 3 -- All puzzles with size > 3 are solvable
   do
      size := n
     create {LINKED_LIST [SOLUTION]} solutions.make
      complete (create {SOLUTION}.make_empty)
   ensure
      solutions_exists : not solutions.is_empty
      complete_solutions: across solutions as s all s.item.row_count = n end
   end
feature \{NONE\} -- Implementation
  complete ( partial: SOLUTION)
      -- Find all complete solutions that extend the partial solution 'partial'
      -- and add them to 'solutions'.
   require
      partial\_exists : partial /= Void
   local
      c: INTEGER
   do
      if partial.row_count = size then
        solutions.extend (partial)
     else
       from
         c := 1
       until
         c > size
       loop
         if not under_attack (partial, c) then
           complete ( partial . extended_with (c))
         end
         c := c + 1
       end
     end
   end
  under_attack (partial: SOLUTION; c: INTEGER): BOOLEAN
      -- Is column 'c' of the current row under attack
      -- by any queen already placed in partial solution 'partial'?
   require
      partial\_exists : partial /= Void
   local
      current_row, row: INTEGER
   do
      current_row := partial.row_count + 1
     from
       row := 1
     until
       Result or row > partial.row_count
     loop
       Result := attack_each_other (row, partial.column_at (row), current_row, c)
```

```
row := row + 1
end
end
attack_each_other (row1, col1, row2, col2: INTEGER): BOOLEAN
-- Do queens in positions ('row1', 'col1') and ('row2', 'col2') attack each other?
do
Result := row1 = row2 or
col1 = col2 or
(row1 - row2).abs = (col1 - col2).abs
end
```

 \mathbf{end}